An Attack Surface Metric

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Context: Security Metrics

Software vendors are spending big on security.



Gauging progress is critical for secure software development. We need measurements and metrics.

We Need Metrics Now!

• A long standing research challenge [ACSAC 01, CRA 03, DIMACS 03, CSTB 07]

Toward a Safer and More Secure Cyberspace [CSTB 2007]:

"..though many benefits would flow from the invention of good metrics, the challenge in this cybersecurity research area is particularly great, and some very new ideas will be needed if cybersecurity metricians are to make more progress."



Measure the system's attack surface

Motivation: ASM is Useful to both Industry and Consumers

A guide in **consumers'** decision making process

A tool in the software development lifecycle to improve security

• design, implementation, testing, deployment, and maintenance

Attack Surface Reduction (ASR) Mitigates Risk

Traditional industry approach: code quality improvement

Software will ship with known and future vulnerabilities

Reduce attack surface to increase the difficulty and decrease the impact of future exploitation

Code Quality and ASR Complement t Each Other

Bad	Medium Security Risk	High Security Risk		
Code				
Quality				
Good	Low Security Risk	Medium Security Risk		
	Low	High		
	Attack Surface Measurement			

Inspiration: Relative Attack Surface Quotient for 7 Versions of Windows [HPW03]



Linux Attack Surface Measurements

Attack Vector		Debian		RH Default		RH Facilities	RH Used
Open socket		15		12		40	41
Open RPC endpoint	[3		3		3	3
Services running as root		21		26		29	30
Services running as nonroot		3		6		8	8
Setuid root programs		54		54		72	72
Local user accounts		21		25		33	34
User id = root accounts		0		4		3	3
Unpassworded accounts		0		0		2	2
Nobody account		1		1		1	1
Weak file permission		7		7		21	37
Scripts enabled				×		2	2

Confirms perception that Debian is more secure than RedHat

Lessons Learned from Windows and Linux Measurements

• Measurement method is ad-hoc

• Requires a security expert

 Focus is on measuring the attack surfaces of operating systems

Research Goals

- Formalize the notion of attack surface
- Introduce a systematic attack surface measurement method
 - Anyone, anywhere, anything
- Validate the method
- Demonstrate the uses of the method

Intuition Behind Attack Surfaces



Hence we define a system's attack surface in terms of the system's resources (i.e., methods, channels, and data items).

Model of a System and its Environment

A system, s, and its environment, $E_s = \langle U, D, T = \{t_1, t_2\} \rangle$.



Formal model uses I/O automata [LT89].

Not All Resources Are Part of the Attack Surface

 Only those resources that the attacker can use to send data into or receive data from the system are relevant.

• We introduce the formal entry point and exit point framework to identify the relevant resources.

Entry Point and Exit Point Framework

- Entry Points/Exit Points
 - Direct (input/output action)
 - Indirect (internal action)



- Channels (e.g., sockets and pipes)
 c ∈ Res(m.pre)



Attack Surface Definition

- Definition
 - M: set of entry points and exit points
 - C: set of channels
 - I: set of untrusted data items.

attack surface = (M, C, I)

Theorem: Given an environment, E, if $AS(A) \ge AS(B)$, then $Attacks(A||E) \supseteq Attacks(B||E)$.

Not All Resources Contribute Equally to the Attack Surface

• Contribution ∞ Damage Potential

Contribution \propto (Attacker Effort) ⁻¹

• Contribution = Damage Potential Attacker Effort

Higher Damage Potential \Rightarrow Stronger m.post \Rightarrow more methods can follow m Lower Attacker Effort \Rightarrow Weaker m.pre \Rightarrow m can follow more methods

Attack Surface Measurement (ASM)

 ASM(A) ≥ ASM(B) if there exists a nonempty set, R, of resources s.t.

 $\forall r \in R. contribution(r, A) \ge contribution(r, B).$

Theorem: Given an environment, E, if $ASM(A) \ge ASM(B)$, then $Attacks(A||E) \supseteq Attacks(B||E)$.

Quantitative Attack Surface Measurement

• Assume der: method \rightarrow Q.

- Similarly, for channel and data.

 $ASM = \left\langle \sum_{m \square M} der(m), \sum_{c \square C} der(c), \sum_{d \square I} der(d) \right\rangle$

Analogous to risk modeling



Abstract Measurement Method

- Identify a set, M, of entry points and exit points, a set, C, of channels, and a set, I, of untrusted data items.
- 2. Estimate each relevant resource's damage potential-effort ratio, der.
- **3.** Compute Attack Surface Measurement =

$$\left\langle \sum_{m \,\square\,M} der(m), \sum_{c \,\square\,C} der(c), \sum_{d \,\square\,I} der(d) \right\rangle.$$

C Measurement Method and Examples

• FTP Servers

- ProFTP 1.2.10 , Wu-FTP 2.6.2

- IMAP Servers
 - Courier 4.0.1, Cyrus 2.2.10

Step 1: Identify Relevant Resources

- Entry Points and Exit Points
 - Static analysis
 - C library methods (e.g., read) for data exchange
 - Call graph
- Channels and Untrusted Data Items
 - Run time monitoring
 - Open channels
 - Data read and written

Step 2: Damage Potential-Effort Ratio

Resource	Damage	Attacker Effort
	Potential	
Method	Privilege	Access Rights
Channel	Protocol	Access Rights
Data Items	Туре	Access Rights

Impose a total ordering among the values of the attributes and assign numeric values accordingly, e.g.,

root = 5 and auth = 3.

FTP Measurement Results

ProFTP = $\langle 312.9, 1.0, 18.9 \rangle$, **Wu-FTP** = $\langle 392.3, 1.0, 17.6 \rangle$



Use domain knowledge to decide which dimension presents more risk and choose accordingly.

Validation

- Validating a software measure is hard [KPF97,....]
 - security metric is even harder

Software	Attack	MS Bulletins, Expert
measure	surface	Survey
Prediction System	Security Risk	IO Automata Model, Patch Analysis, Anecdotal Evidence

Liu and Traore independently validated our metric [LT07].

Validating the Measurement Method

Key Assumptions

- Three dimensions of the attack surface
- Damage potential-effort ratio
- Six attributes
 - method privilege, method access rights, channel protocol, channel access rights, data item type, and data item access rights

Statistical Analysis of Microsoft Security Bulletins (MSB)

- An MSB mentions a vulnerability and resources needed for exploitation
- Are methods, channels, and data used in the exploitation?
- Analyzed MSBs from 2004-2006

Methods	
Channels	
Data	

Results: The Attributes are Indicators of Damage Potential and Effort

Attribute	Significance	Correlation
Privilege		
Method Access Rights		
Protocol		?
Channel Access Rights		
Туре		?
Data Access Rights		

Expert Linux System Administrator Survey

- MSB has no data relevant to a resource's attackability
 - Could not validate damage potential-effort ratio
- Surveys are widely used to collect a wide range of data
 - Prior work uses surveys to validate measures [K87,]
 - Feedback from one target user group (Industrial collaboration for other target user group)
 - W.r.t. Linux (MSB w.r.t. Windows)

Results: A Majority of the Subjects Agree With Our Measurement Method

Methods		Privilege	~
Channels		Method Access Rights	<
Data		Protocol	?
Damage		Channel Access Rights	\checkmark
Potential-		Туре	?
Effort Ratio		Data Access Rights	\checkmark

Validating the Prediction System

- Show that if system A is more secure than system B, then ASM(A) < ASM(B)
- Assumption: Vulnerability patches improve software security
 - ASM(After Patch) < ASM(Before Patch)</p>

Patches reduce attack surface measurement

Results: A Majority of the Patches Reduce ASM

Software	Percentage of Patches that reduce ASM	Significance (p< 0.05)
Firefox 2.0	67%	\checkmark
ProFTP (all)	70%	\checkmark
All NVD Bulletins	76.9%	\checkmark

Anecdotal Evidence from Industry

- Microsoft
 - Sasser Worm
 - Nachi Worm
 - Zotob Worm
- Firefox 2.0
 - SSL buffer overflow

Collaboration with SAP

- SAP is world's largest provider of enterprisescale software
 - Complex technology platforms and business applications
- Demonstrate that the measurement method scales to enterprise-scale software
- Receive feedback from software architects and developers

Java Measurement Tool Screenshot

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Results

- Measured the attack surface of a key component of SAP component
 - Measurement results conform to expectation
 - Detailed tool output, incremental analysis, and what-if scenarios are useful for attack surface reduction
 - Lessons learned

ASM in Software Development LifeCycle

Im

Compare and reduce ASM from version to version [Microsoft, Firefox, OpenSSH] Use ASM to guide testing and code inspection [MuSecurity, SAP]

irefox, Use ASM to choose ^{-]} a secure **configuration** Use ASM in patch [⊂]irefox] implementation

Future Work: Software Development

• Range analysis



- Other uses
 - ``Safe'' software composition
 - Testing, deployment, maintenance

Future Work: Software Consumers

- Attack surface measurement in the absence of source code
 - Components as Entry/Exit points
 - Channels and Data as before
- Multiple metrics are needed for decision support

How do we combine multiple measures?

Related Work-1

 Prior work assumes the knowledge of vulnerabilities [AB95, VGMCM96, ODM99...]

- ASM is based on a system's inherent properties
 - Formal framework encompasses past, present, and future vulnerabilities
 - Complementary to prior work

Related Work-2

 Prior work takes an attacker-centric approach [S99, MBFB05, LB08,..]

- ASM takes a system-centric approach
 - Depends on a system's design
 - No assumptions about the attacker
 - Can be used as a tool in software development

Related Work-3

 Prior work is conceptual in nature and haven't been applied to real systems [AB95, MGVT02, S04,..]

- We measured the attack surfaces of real-world software
 - FTP servers, IMAP servers
 - SAP business applications

Summary

- Introduced a pragmatic approach for security measurement
 - Software industry found it useful [Microsoft, Firefox, OpenSSH, MuSecurity, SAP, ..]
- First step in the grander challenge of security metrics
 - Understanding over time will lead to more meaningful metrics

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Backups

I/O Automata [LT89]

- Action Signature
 - Input, Output, Internal actions
 - Pre and Post conditions
 m.pre and m.post



• Composition

$$- E_{s} = (U_{io} || D_{io} || (||t_{io})) - P = s_{io} || E_{s} t_{io} T_{io}$$

Validation of the Attributes

• An MSB has a severity rating and mentions the six resources attributes

Impact	Damage Potential
Difficulty	Attacker Effort

Significant Predictor	Two sided Z-test (p < 0.05)
Correlation	Sign of Coefficient in Ordered Logistic Regression

Inspiration: Howard's Relative Attack Surface Quotient (RASQ)[H03]

- Howard's informal RASQ Measurement Method
 - Identify a system's attack vectors
 - Assign weights to the attack vectors to reflect their attackability
 - RASQ = sum of the weighted counts of the attack vectors

Direct Entry Points

Methods that directly receive data.



direct entry point: an input action with a matching output action

Indirect Entry Points

Methods that indirectly receive data.



indirect entry point: internal action (m1.post => m.post) Λ (d ε Res(m1.post) Λ d ε Res(m.pre))

Channels and Data

Channels (e.g., sockets and pipes) • c ext{c} Res(m.pre)





Definition of An Attack

Attacks $(s_{io}) =$ Set of executions of $(s_{io} || E_s)$ that contain either an input action or output action of s_{io} .

Not All Resources Contribute Equally to the Attack Surface

- contribution α damage potential α 1/attacker effort
- r1 ≥ r2 if higher damage potential and/or lower attacker effort

m(MA, CA, DA, MB, CB, DB)

pre: $P_{pre} \wedge (MA \ge m.ef) \wedge (CA \ge c.ef) \wedge (DA \ge d.ef)$

post: $P_{post} \land (MB \ge m.dp) \land (CB \ge c.dp) \land (DB \ge d.dp)$

Damage Potential-Effort Ratio

• Contribution \propto Damage Potential Contribution \propto (Attacker Effort) ⁻¹



C Measurement Method



Survey Methodology

- Email survey of experienced Linux system administrators
 - Diverse background and geographic location
- Questions on a five point Likert scale [L32]
 - Pretesting and interviewing to avoid bias
 - Self-selection bias
- **Descriptive** analysis techniques
 - Central tendency bias
 - %age of agreement, disagreement, and neither
 - t-test (p < 0.05)</p>

Not All Patches Are Relevant

- Heuristics: vulnerability type determines patch relevance
 - Use National Vulnerability Database (NVD) type information
 - Infer type if missing
- Not all relevant patches reduce the attack surface
- Consider local effect of a patch

Data Collection for Firefox 2.0



Java Measurement Method

- Focus on method dimension
- Entry Points and Exit Points
 - Call graph
 - Interface methods, methods invoking other systems' interfaces and Java I/O library methods
- Damage Potential-Effort Ratio
 - Use SAP's threat modeling process to assign numbers

Tool Usage in Software Development

- Tool produces detailed output
 - Guides attack surface reduction

Incremental analysis

- What-If scenarios
 - Addition of a new feature
 - Removal of a feature



FTP Daemons (method)

- 1. Access rights don't matter.
- 2. proftpd privilege level contributes more than auth.'s.



Tool Output

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	Measurement
	Attack Surface Measurement
Contribution	Weights assigned to attributes: Parameter: 35.0 Data Store Invocation: 18.0 Other System Invocation: 1.0 Access Rights: 1.0 Total number of source methods in the call graph: 670 Datal number of entry points and exit points: 71
E	Total Attack Surface Measurement: 5279.0 List of Entry Points and Exit Points Fully Qualified Method Name (parameter, data store, other systems) ASM Contribution (3, 0, 0) 105.0
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